A Proximity Algorithm Solving Indicator Functions Based $\ell_1$-Norm Minimization Problems in Compressive Sampling

Feishe Chen†  Lixin Shen†  Bruce W. Suter‡  Yuesheng Xu†

Abstract

An accurate and efficient algorithm for an $\ell_1$-norm minimization problem is highly needed and is crucial for the success of sparse signal recovery in compressive sampling, a recent development in the field of data analysis. Most of existing algorithms in the literature give an approximated solution to the problem. We tackle the $\ell_1$-norm minimization problem by equivalently reformulating it via an indicator function which describes the constraints for the problem. It turns out that the resulting model can be solved efficiently and accurately by using an elegant proximity operator based algorithm. We establish the convergence analysis of the resulting algorithm. Numerical experiments show that the proposed algorithm performs well for sparse signals with magnitudes over a high dynamic range. Furthermore, it performs significantly better than the well-known algorithm NESTA in terms of the quality of restored signals and the computational complexity measured in the CPU-time consumed.

1 Introduction

In this paper, we study recovery of an unknown vector $u_0 \in \mathbb{R}^n$ from the observed data $b \in \mathbb{R}^m$ and the model

$$b = Au_0 + z,$$

where $A$ is a known $m \times n$ measurement matrix and $z \in \mathbb{R}^m$ is a noise term. Under an assumption that the vector $u_0$ of interest is sparse, the work in [3, 4, 5] shows that an accurate estimation of $u_0$ is possible even when $m < n$, that is, the observations are less than unknowns. Recently, there is a significant body of work that focuses on finding an approximation to $u_0$ by solving a convex optimization problem. In the presence of noise-free data, i.e., $z = 0$, this optimization problem is

$$(BP) \quad \min \{ \| u \|_1 : u \in \mathbb{R}^n \} \quad \text{subject to} \quad b = Au,$$

which essentially is the basis pursuit problem proposed early in the context of time-frequency representation [8]. Here, $\| \cdot \|_1$ denotes the $\ell_1$-norm of a vector in an Euclidean space. The optimization model (BP) can be solved by linear programming.

In the presence of noisy data, the linear constraint $b = Au$ in (BP) is relaxed to an inequality constraint $\|Au - b\|_2 \leq \epsilon$, where $\| \cdot \|_2$ denotes the $\ell_2$-norm of a vector in an Euclidean space. As a result, the optimization model (BP) becomes the basis pursuit denoising problem

$$(BP_{\epsilon}) \quad \min \{ \| u \|_1 : u \in \mathbb{R}^n \} \quad \text{subject to} \quad \|Au - b\|_2 \leq \epsilon,$$

$\ast$This research is supported in part by 2012 Air Force Visiting Faculty Research Program, by an award from the NRC (National Research Council) via AFSOR, by the US National Science Foundation under grant DMS-1115523.

†Department of Mathematics, Syracuse University, Syracuse, NY 13244, USA.

‡Air Force Research Laboratory, AFRL/RITB, Rome, NY 13441-4505. Email: Bruce.Suter@rl.af.mil
where $\epsilon^2$ is an estimated upper bound of the noise power. Apparently, $(BP_\epsilon)$ reduces to $(BP)$ when $\epsilon = 0$.

Both problems $(BP)$ and $(BP_\epsilon)$ are closely related to the penalized least-squares problem

$$(QP_\lambda) \quad \min \left\{ \frac{1}{2} \|Au - b\|_2^2 + \lambda \|u\|_1 : u \in \mathbb{R}^n \right\}.$$ 

It is well-known that a solution of $(QP_\lambda)$ converges to a solution of $(BP)$ when $\lambda$ approaches to 0. Furthermore, it is stated in convex analysis (see, e.g., [23]) that a solution of $(BP_\epsilon)$ is either 0, or a minimizer of $(QP_\lambda)$ for some $\lambda > 0$ depending on $\epsilon$.

A large amount of research has been done on solving problems $(BP)$, $(BP_\epsilon)$, and $(QP_\lambda)$. Here, we only give a brief and non-exhaustive review of results for these problems. In [8], problems $(BP)$ and $(QP_\lambda)$ are solved by first reformulating them as perturbed linear programs and then applying a primal-dual interior-point approach [26]. Recently, many iterative shrinkage/thresholding algorithms are proposed to handle problem $(QP_\lambda)$. These include the proximal forward-backward splitting [9], the gradient projection for sparse reconstruction [13], the FISTA (fast iterative shrinkage-thresholding algorithm) [1], the fixed-point continuation algorithm [14], the Bregman iterative regularization [27], and the reference therein. Problem $(BP_\epsilon)$ also frequently appears in wavelet-based signal/image restoration [6, 7, 10] with the matrix $A$ associated with some inverse transforms.

Problem $(BP_\epsilon)$ can be formulated as a second-order cone program and solved by interior-point algorithms. Many suggested algorithms for $(BP_\epsilon)$ are based on repeatedly solving $(QP_\lambda)$ for various values of $\lambda$. Such algorithms are referred to as the homotopy method originally proposed in [12, 22]. The homotopy method is also successfully applied to $(BP)$ in [11]. A common approach for obtaining approximate solutions to $(BP_\epsilon)$ is often accompanied by solving $(QP_\lambda)$ for a decreasing sequence of values of $\lambda$ [24]. The optimization theory asserts that problems $(BP_\epsilon)$ and $(QP_\lambda)$ are equivalent provided that the parameters $\epsilon$ and $\lambda$ satisfy certain relationship [23]. Since this relationship is hard to compute in general, solving problem $(BP_\epsilon)$ via repeatedly solving $(QP_\lambda)$ for various values of $\lambda$ is problematic. Recently, the NESTA [2] which employs Nesterov’s optimal gradient method was proposed for solving relaxed versions of $(BP)$ and $(BP_\epsilon)$ via the Nesterov’s smoothing technique [20]. Clearly, the closeness of the solution to the relaxed version of $(BP)$ (or the relaxed version of $(BP_\epsilon)$) to the solution to $(BP)$ (or $(BP_\epsilon)$) is determined by the level of the closeness of the smoothed $\ell_1$-norm to the $\ell_1$-norm itself. Certainly, the performance of these approaches depends on the fine tuning of the parameter $\lambda$ or a parameter that controls the degree of the closeness of the $\ell_1$-norm and its smoothed version.

In this paper, we propose a new algorithm for solving problems $(BP)$ and $(BP_\epsilon)$. To this end, we convert each constrained optimization problem to an unconstrained one via a indicator function. The corresponding objective function for the unconstrained optimization problem is the sum of the $\ell_1$-norm of the underlying signal $u$ and the indicator function of a set in $\mathbb{R}^m$, which is $\{0\}$ for $(BP)$ or the $\epsilon$-ball for $(BP_\epsilon)$, composing with the affine transformation $Au - b$. Non-differentiability of both the $\ell_1$-norm and the indicator of the set imposes challenges for solving the associated optimization problem. Fortunately enough, their proximity operators have explicit expressions. The solutions for the problem can be viewed as fixed-points of a coupled equation formed in terms of these proximity operators. An iterative algorithm for finding the fixed-points is then developed. The main advantage of this algorithm is that solving $(QP_\lambda)$ or smoothing the $\ell_1$-norm are no longer necessary. This makes the proposed algorithm attractive for solving $(BP)$ and $(BP_\epsilon)$. The efficiency of proximity algorithms has been demonstrated in [9, 15, 16, 17] for various image processing models.

The rest of the paper is organized as follows: In section 2 we reformulate the $\ell_1$-norm minimization problems $(BP)$ and $(BP_\epsilon)$ via an indicator function and characterize solutions of the proposed
model in terms of fixed-point equations. We also point out the connection between the proposed model and (QP_\lambda) in terms of the Moreau envelope. In section 3 we develop an algorithm for the resulting new minimization problem based on the fixed-point equations arising from the characterization of the proposed model. Numerical experiments are presented in section 4. We draw our conclusions in section 5.

2 An \ell_1-Norm Optimization Model via an Indicator Function

In this section, we consider general optimization models that include models (BP) and (BP_\epsilon) as their special cases and characterize solutions to the proposed models.

We begin with introducing our notation and recalling necessary background from convex analysis. For the usual d-dimensional Euclidean space denoted by \( \mathbb{R}^d \) we define \( \langle x, y \rangle := \sum_{i=1}^{d} x_i y_i \), for \( x, y \in \mathbb{R}^d \), the standard inner product of \( \mathbb{R}^d \). The class of all lower semicontinuous convex functions \( f: \mathbb{R}^d \to (-\infty, +\infty] \) such that \( \text{dom} f := \{ x \in \mathbb{R}^d : f(x) < +\infty \} \neq \emptyset \) is denoted by \( \Gamma_0(\mathbb{R}^d) \). The indicator function of a closed convex set \( C \) in \( \mathbb{R}^d \) is defined, at \( u \in \mathbb{R}^d \), as \( \iota_C(u) := \begin{cases} 0, & \text{if } u \in C, \\ +\infty, & \text{otherwise}. \end{cases} \)

Clearly, the indicator function \( \iota_C \) is in \( \Gamma_0(\mathbb{R}^d) \) for any closed nonempty convex set \( C \). In particular, we define a ball in \( \mathbb{R}^m \) centered at the origin with radius \( \epsilon \) as \( B_\epsilon := \{ v : v \in \mathbb{R}^m \text{ and } \| v \|_2 \leq \epsilon \} \).

Given a matrix \( A \in \mathbb{R}^{m \times n} \) and a vector \( b \in \mathbb{R}^m \), we consider the following optimization problem

\[
\min \{ \| u \|_1 + \iota_{B_\epsilon}(Au - b) : u \in \mathbb{R}^n \}. \tag{2}
\]

We can easily see that if \( \epsilon = 0 \) then model (2) reduces to (BP); if \( \epsilon > 0 \) then model (2) reduces to (BP_\epsilon). In other words, the both constrained optimization problems (BP) and (BP_\epsilon) can be unified as the unconstrained optimization problem (2) via the indicator function \( \iota_{B_\epsilon} \).

In the following, we shall focus on characterizing solutions of model (2) using fixed-point equations. To characterize solutions of model (2), we first need two concepts in convex analysis, namely, the proximity operator and subdifferential for functions in \( \Gamma_0(\mathbb{R}^d) \). The proximity operator was introduced in \([18, 19]\). For a function \( f \in \Gamma_0(\mathbb{R}^d) \) at a given vector \( u \in \mathbb{R}^d \), the proximity operator of \( f \) with parameter \( \lambda \), denoted by \( \text{prox}_{\lambda f} \), is a mapping from \( \mathbb{R}^d \) to itself, defined for a given point \( x \in \mathbb{R}^d \) by

\[
\text{prox}_{\lambda f}(x) := \arg\min \left\{ \frac{1}{2\lambda} \| u - x \|_2^2 + f(u) : u \in \mathbb{R}^d \right\}.
\]

The subdifferential of a proper convex function \( \psi \in \Gamma_0(\mathbb{R}^d) \) at a given vector \( u \in \mathbb{R}^d \) is the set defined by

\[
\partial \psi(u) := \{ v : v \in \mathbb{R}^d \text{ and } \psi(w) \geq \psi(u) + \langle v, w - u \rangle \text{ for all } w \in \mathbb{R}^d \}.
\]

The subdifferential and the proximity operator of the function \( \psi \) are related in the following way (see, e.g. \([16]\)): for \( u \) in the domain of \( \psi \) and \( v \in \mathbb{R}^d \)

\[
v \in \partial \psi(u) \text{ if and only if } u = \text{prox}_\psi(u + v). \tag{4}
\]

Now, with the help of the subdifferential and the proximity operator, we can characterize a solution of the indicator function based model (2) via fixed-point equations.
Proposition 2.1 Let $\epsilon$ be a nonnegative number, let $B_\epsilon$ be the ball in $\mathbb{R}^m$ centered at the origin with radius $\epsilon$, let $b$ be a point in $\mathbb{R}^m$, and let $A$ be an $m \times n$ matrix. If $u \in \mathbb{R}^n$ is a solution to model (2), then for any $\alpha > 0$ and $\beta > 0$ there exists a vector $v \in \mathbb{R}^m$ such that

$$u = \text{prox}_{\frac{\beta}{\alpha} \cdot \| \cdot \|_1} \left( u - \frac{\beta}{\alpha} A^\top v \right), \quad (5)$$

$$v = (I - \text{prox}_{\frac{1}{\alpha} \cdot \| \cdot \|_1})(Au + v). \quad (6)$$

Conversely, if there exist $\alpha > 0$, $u \in \mathbb{R}^n$, and $v \in \mathbb{R}^m$ satisfying equations (5) and (6), then $u$ is a solution of model (2).

Proof: We first assume that $u \in \mathbb{R}^n$ is a solution to model (2). Set $\varphi := \iota_{B_\epsilon} \cdot (-b)$. Hence $Au - b$ must be in the ball $B_\epsilon$. Therefore, both sets $\partial \| \cdot \|_1(u)$ and $\partial \varphi(Au)$ are nonempty. By Fermat’s rule we have that $0 \in \partial \| \cdot \|_1(u) + A^\top \partial \varphi(Au)$. Therefore, for any $\alpha > 0$ and $\beta > 0$ there exist $w \in \frac{1}{\alpha} \partial \| \cdot \|_1(u)$ and $v \in \frac{1}{\beta} \partial \varphi(Au)$ such that $0 = \alpha w + \beta A^\top v$, i.e., $w = -\frac{\beta}{\alpha} A^\top v$. By using (4), inclusion $w \in \frac{1}{\alpha} \partial \| \cdot \|_1(u)$ implies $u = \text{prox}_{\frac{1}{\alpha} \cdot \| \cdot \|_1}(u + w)$, which is equation (5). Since $\frac{1}{\beta} \varphi = \varphi$ for any $\beta > 0$, inclusion $v \in \frac{1}{\beta} \partial \varphi(Au)$ leads to $Au = \text{prox}_{\frac{1}{\beta} \cdot \| \cdot \|_1}(v + Au)$, which is equivalent to (6).

Conversely, if equations (5) and (6) are satisfied for some $\alpha > 0$, $\beta > 0$, $u \in \mathbb{R}^n$, and $v \in \mathbb{R}^m$, using (4) again, we have that $-\frac{\beta}{\alpha} A^\top v \in \partial \left( \frac{1}{\alpha} \cdot \| \cdot \|_1 \right)(u)$ and $v \in \partial \varphi(Au)$. Since $\partial \left( \frac{1}{\alpha} \cdot \| \cdot \|_1 \right)(u) = \frac{1}{\alpha} \partial \| \cdot \|_1(u)$ and $\partial \varphi(Au) = \beta \partial \varphi(Au)$, we know from the above that $0 \in \partial \| \cdot \|_1(u) + A^\top \partial \varphi(Au)$. This indicates that $u$ is a solution of model (2). The proof is complete.

We remark that the above fixed-point characterization can be identified as a special case of Proposition 1 in [17]. We include the proof of Proposition 2.1 here for making the paper self-contained.

The characterization involves the proximity operators of the functions $\| \cdot \|_1$ and $\iota_{B_\epsilon} \cdot (-b)$. To make this characterization useful, it requires that both proximity operators should be computed easily and efficiently. Indeed, the proximity operator $\text{prox}_{\frac{1}{\alpha} \cdot \| \cdot \|_1}$ is the soft-thresholding operator in wavelet analysis and is given at $u \in \mathbb{R}^n$ as follows:

$$\left( \text{prox}_{\frac{\beta}{\alpha} \cdot \| \cdot \|_1}(u) \right) [i] = \max \left\{ |u[i]| - \frac{1}{\alpha} \epsilon, 0 \right\} \text{sign}(u[i]), \quad \text{for } i = 1, 2, \ldots, n. \quad (7)$$

The proximity operator $\text{prox}_{\iota_{B_\epsilon} \cdot (-b)}$ is given by the following lemma.

Lemma 2.2 Let $\epsilon$ be a nonnegative number, let $B_\epsilon$ be the ball in $\mathbb{R}^m$ centered at the origin with radius $\epsilon$, let $b$ be a point in $\mathbb{R}^m$. Then for a given $v \in \mathbb{R}^m$

$$\text{prox}_{\iota_{B_\epsilon} \cdot (-b)}(v) = b + \min \left\{ 1, \frac{\epsilon}{\| v - b \|_2} \right\} \cdot (v - b).$$

Proof: By the definition of the proximity operator, we have that for a vector $v \in \mathbb{R}^m$

$$\text{prox}_{\iota_{B_\epsilon} \cdot (-b)}(v) = \arg\min \left\{ \frac{1}{2} \| y - v \|_2 + \iota_{B_\epsilon}(y - b) : y \in \mathbb{R}^m \right\} = b + \arg\min \left\{ \frac{1}{2} \| \tilde{y} - (v - b) \|_2 + \iota_{B_\epsilon}(\tilde{y}) : \tilde{y} \in \mathbb{R}^m \right\}.$$

The second term of the last equation of the above is the projection of the vector $v - b$ onto the ball $B_\epsilon$. Hence, that is equal to $\min \left\{ 1, \frac{\epsilon}{\| v - b \|_2} \right\} \cdot (v - b)$. The proof is complete. \qed
We close this section by making a connection between the proposed model (2) and model (QP_{\lambda}) via the Moreau envelope [18]. Recall that the Moreau envelope of a function \( f \in \Gamma_0(\mathbb{R}^d) \) with parameter \( \mu \) at \( x \in \mathbb{R}^d \) is
\[
\text{env}_\mu f(x) := \min \left\{ \frac{1}{2\mu} \| y - x \|^2_2 + f(y) : y \in \mathbb{R}^d \right\},
\]
which is also in \( \Gamma_0(\mathbb{R}^d) \). In particular, the Moreau envelope of the indicator function \( \iota_{B_{\epsilon}} \) with parameter \( \mu \) is given in the following result.

**Lemma 2.3** Let \( \epsilon \) be a nonnegative number, let \( \mu \) be a positive number, let \( B_{\epsilon} \) be the ball in \( \mathbb{R}^m \) centered at the origin with radius \( \epsilon \). Then for any \( x \in \mathbb{R}^m \)
\[
\text{env}_\mu \iota_{B_{\epsilon}}(x) = \frac{1}{2\mu} \left( \max\{\|x\|_2 - \epsilon, 0\} \right)^2.
\]

**Proof:** By the definition of the Moreau envelope, we have that
\[
\text{env}_\mu \iota_{B_{\epsilon}}(x) = \min \left\{ \frac{1}{2\mu} \| y - x \|^2_2 + \iota_{B_{\epsilon}}(y) : y \in \mathbb{R}^m \right\}.
\]
Hence, \( \text{env}_\mu \iota_{B_{\epsilon}}(x) \) is the square of the distance from the point \( x \) to the ball \( B_{\epsilon} \) then scaled by the factor \( \frac{1}{2\mu} \). Since this distance is either 0 if \( x \) is inside the ball or \( \|x\|_2 - \epsilon \) if \( x \) is outside the ball, our result follows immediately. \( \square \)

![Figure 2.1: The plots of env_{\mu \iota_{\{0\}}}, the Moreau envelope of the indicator function \iota_{\{0\}} with \mu = 8, 1, 1/8.](image)

It is not difficult to see from Lemma 2.3 that the envelope \( \text{env}_{\mu \iota_{B_{\epsilon}}}(x) \) converges pointwisely to the indicator function \( \iota_{B_{\epsilon}}(x) \) when \( \mu \) approaches to 0 for every point \( x \) as depicted in Figure 2.1 for \( \epsilon = 0 \) and Figure 2.2 for some \( \epsilon > 0 \). In particular, we have from Lemma 2.3 that
\[
\text{env}_{\lambda \iota_{\{0\}}} = \frac{1}{2\lambda} \| \cdot \|^2_2.
\]
Hence, problem (QP$_\lambda$) can be rewritten in terms of the Moreau envelope as
\[
\min \left\{ \|u\|_1 + \text{env}_{\lambda \iota_{(0)}}(Au - b) : u \in \mathbb{R}^n \right\}.
\] (8)

Surprisingly, model (8) (i.e., (QP$_\lambda$)) is used for finding an approximate solution to both problems (BP) and (BP$_\epsilon$) in the literature. Note that
\[
\iota_{B_\epsilon}(x) - \text{env}_{\lambda \iota_{(0)}}(x) = \iota_{B_\epsilon}(x)
\]
and
\[
\iota_{B_\epsilon}(x) - \text{env}_{\lambda \iota_{(0)}}(x) = \begin{cases} 
-\frac{1}{2\lambda} \|x\|_2^2, & \text{if } x \in B_\epsilon; \\
+\infty, & \text{otherwise}.
\end{cases}
\]

These identities indicate that env$_{\lambda \iota_{B_\epsilon}}$ approximates $\iota_{B_\epsilon}$ clearly better than env$_{\lambda \iota_{(0)}}$. Therefore, we comment that instead of model (8) the following model
\[
\min \left\{ \|u\|_1 + \text{env}_{\lambda \iota_{B_\epsilon}}(Au - b) : u \in \mathbb{R}^n \right\}
\] (9)
would be more suitable for finding an approximate solution to problem (BP$_\epsilon$). In other words, (QP$_\lambda$) used in many algorithms (e.g., [12, 22, 24, 25]) for (BP$_\epsilon$) should be replaced by model (9). By doing so, an improvement for solving (BP$_\epsilon$) should be expected if the noise power $\epsilon^2$ can be estimated in advance. This will be studied in our future project. Unlike models (8) and (9) we solve model (2) directly without introducing the parameter $\lambda$.

3 An Algorithm and Its Convergence

In this section, we develop an algorithm for finding a solution of model (2) and give the convergence analysis of the developed algorithm.

As we already know, all solutions of model (2) should satisfy the fixed-point equations given by (5) and (6). Our proposed algorithm for model (2) is derived based on the following equivalent
form of equations (5) and (6)

\[
\begin{align*}
    u &= \prox_{\frac{1}{\alpha} \| \cdot \|_1} \left( \left( I - \frac{\beta}{\alpha} A^\top A \right) u - \frac{\beta}{\alpha} A^\top (v - w) \right), \\
    w &= \prox_{\beta \| \cdot \|_2} (Au + v), \\
    v &= Au + v - w.
\end{align*}
\]

(10)

Based on the above fixed-point equations in terms of \( u, w, \) and \( v, \) for arbitrary initial vectors \( u^0 \in \mathbb{R}^n, w^0, v^0 \in \mathbb{R}^m, \) we generate the sequence \( \{u^k : k \in \mathbb{N}_0\} \) by the following iterative scheme

\[
\begin{align*}
    u^{k+1} &= \prox_{\frac{1}{\alpha} \| \cdot \|_1} \left( \left( I - \frac{\beta}{\alpha} A^\top A \right) u^k - \frac{\beta}{\alpha} A^\top (v^k - w^k) \right), \\
    w^{k+1} &= \prox_{\beta \| \cdot \|_2} (Au^{k+1} + v^k), \\
    v^{k+1} &= Au^{k+1} + v^k - w^{k+1},
\end{align*}
\]

(11)

where \( \mathbb{N}_0 := \{0, 1, \ldots \}. \)

To show the convergence of the iterative scheme (11), we recall a result from a recent paper [15].

**Theorem 3.1 (Theorem 3.5 in [15])** If \( x \) is a vector in \( \mathbb{R}^n, \) \( A \) is an \( m \times n \) matrix, \( \varphi \) is a proper convex function on \( \mathbb{R}^m \) and \( \alpha, \beta, \lambda \) are positive numbers such that

\[
\frac{\beta}{\lambda \alpha} < \frac{1}{\| A \|^2},
\]

(12)

then the sequence \( \{u^k : k \in \mathbb{N}_0\} \) generated by the following iterative scheme

\[
\begin{align*}
    u^{k+1} &= x + \prox_{\frac{1}{\lambda} \| \cdot \|_1} \left( \left( I - \frac{\beta}{\lambda \alpha} A^\top A \right) u^k - x - \frac{\beta}{\lambda \alpha} A^\top (v^k - w^k) \right), \\
    w^{k+1} &= \prox_{\beta \| \cdot \|_2} (Au^{k+1} + v^k), \\
    v^{k+1} &= Au^{k+1} + v^k - w^{k+1}
\end{align*}
\]

(13)

converges to a solution of the optimization problem

\[
\min \{\lambda \| u - x \|_1 + \varphi \circ A(u) : u \in \mathbb{R}^m\}.
\]

(14)

With the help of Theorem 3.1, the following result shows that under appropriate conditions on parameters \( \alpha \) and \( \beta \) the sequence \( \{u^k : k \in \mathbb{N}_0\} \) converges to a solution of model (2).

**Theorem 3.2** Let \( \epsilon \) be a nonnegative number, let \( B_\epsilon \) be the ball in \( \mathbb{R}^m \) centered at the origin with radius \( \epsilon, \) let \( b \) be a point in \( \mathbb{R}^m, \) and let \( A \) be an \( m \times n \) matrix. If

\[
\frac{\beta}{\alpha} < \frac{1}{\| A \|^2},
\]

(15)

then for arbitrary initial vectors \( u^0 \in \mathbb{R}^n, w^0, v^0 \in \mathbb{R}^m, \) the sequence \( \{u^k : k \in \mathbb{N}_0\} \) generated from the iterative scheme (11) converges to a solution of model (2).

**Proof:** By setting \( x = 0 \) and \( \lambda = 1 \) and identifying \( \varphi = \mathbb{1}_{B_\epsilon} (\cdot - b) \) in model (14), the iterative scheme (11) can be viewed as a special case of the one given in (13). Our result follows immediately from Theorem 3.1. \( \square \)

The convergence result given by Theorem 3.2 offers a practical way to find a solution of model (2). Since the explicit forms of the proximity operators \( \prox_{\frac{1}{\alpha} \| \cdot \|_1} \) and \( \prox_{\beta \| \cdot \|_2} \) are given
Algorithm 1 (An iterative scheme for model (BP) with $\epsilon \geq 0$)

Initialization: $v^0 \in \mathbb{R}^m$, $w^0 \in \mathbb{R}^m$, $u^0 \in \mathbb{R}^n$, $\epsilon > 0$, $\alpha > 0$, and $\beta > 0$ with $\frac{\beta}{\alpha} < \frac{1}{\|A\|^2}$.

repeat ($k \geq 0$)

Step 1: Compute

$$u^{k+1} = \text{prox}_{\frac{1}{\alpha} \| \cdot \|_1} \left( \left( I - \frac{\beta}{\alpha} A^\top A \right) u^k - \frac{\beta}{\alpha} A^\top (v^k - w^k) \right)$$

Step 2: Compute

$$w^{k+1} = \begin{cases} Au^{k+1} + v^k, & \text{if } \|Au^{k+1} + v^k - b\|_2 < \epsilon; \\ b + \epsilon \frac{Au^{k+1} + v^k - b}{\|Au^{k+1} + v^k - b\|_2}, & \text{otherwise.} \end{cases}$$

Step 3: Compute

$$v^{k+1} = Au^{k+1} + v^k - w^{k+1}$$

until a given stopping criteria is met

by (7) and Lemma 2.2, respectively, based on Theorem 3.2 a unified approach for solving both (BP) and (BP) is depicted in Algorithm 1.

We make some comments on Algorithm 1. Step 1 of computing $u^{k+1}$ is from the first equation in (11); Step 2 of computing $w^{k+1}$ is from the second equation in (11) and Lemma 2.2; Step 3 of computing $v^{k+1}$ is exactly the same as the last equation in (11). This algorithm can be presented in a more computationally efficient way by combining Step 2 and Step 3 together and eliminating the intermediate variable $w^k$. The motivation comes from the observation $Au^k - w^k = v^k - v^{k-1}$ which is due to the third step of Algorithm 1. Substituting $Au^k - w^k$ in Step 1 by $v^k - v^{k-1}$ yields that

$$u^{k+1} = \text{prox}_{\frac{1}{\alpha} \| \cdot \|_1} \left( u^k - \frac{\beta}{\alpha} A^\top (2v^k - v^{k-1}) \right),$$

(16)

with an assumption $v^{-1} = v^0 - (Au^0 - w^0)$ for given $u^0$, $w^0$, and $v^0$. We, further, can substitute $w^{k+1}$ computed in Step 2 into Step 3. In such the way, the intermediate variable $w^k$ is no longer needed. Hence, these simplification yield Algorithm 2, a variant of Algorithm 1. When $\epsilon = 0$, all vectors $w^{k+1}$ in Algorithm 1 are equal to the constant vector $b$ for all $k \geq 0$. Because of this, we would like to set $w^0 = b$ in both Algorithms 1 and 2. Finally, it is more efficient to update $u^{k+1}$ with Step 1 of Algorithm 2 than Step 1 of Algorithm 1 in each iteration since the matrix-vector multiplication involving $A$ is not required in equation (16). However, updating $u^{k+1}$ via the formulation of Step 2 in Algorithm 1 can be implemented through the use of the component-wise Gauss-Seidel iteration which may accelerate the rate of convergence of the algorithm and therefore reduce the total CPU-time consumed. The efficiency of component-wise Gauss-Seidel iteration has been verified in [15, 16].

To use Algorithm 2, one needs to fix the parameters $\alpha$ and $\beta$ such that $\frac{\beta}{\alpha} < \frac{1}{\|A\|^2}$ (see Theorem 3.2). From Step 1 of the algorithm, the ratio $\frac{\beta}{\alpha}$ plays a role of step-size of changing $u^k$. Therefore, we keep this ratio as big as possible and set

$$\beta = 0.999 \|A\|^2 \alpha$$

(17)

in our numerical experiments. In such the way, $\alpha$ is essentially the only parameter that needs to be determined. We now investigate the impact of the parameter $\alpha$ on the performance of Algorithm 2.
Algorithm 2 (A variant of Algorithm 1 for model (BP$_{\epsilon}$))

Initialization: $v^0 \in \mathbb{R}^m$, $u^0 \in \mathbb{R}^n$, $\epsilon > 0$, $\alpha > 0$, and $\beta > 0$ with $\frac{\beta}{\alpha} < \frac{1}{\|A\|^2}$; set $v^{-1} = v^0 - (Au^0 - d^0)$,

repeat ($k \geq 0$)

Step 1: Compute $u^{k+1} = \text{prox}_{\frac{1}{\alpha\|\cdot\|_1}}\left(u^k - \frac{\beta}{\alpha} A^\top (2v^k - v^{k-1})\right)$

Step 2: Compute $v^{k+1} = \begin{cases} 0, & \text{if } \|Au^{k+1} + v^{k+1} - b\|_2 < \epsilon; \\ \left(1 - \frac{\|Au^{k+1} + v^{k+1} - b\|_2}{\|Au^k + v^k - b\|_2}\right) (Au^{k+1} + v^k - b), & \text{otherwise}. \end{cases}$

until a given stopping criteria is met

For this purpose, we set up the configurations of our experiments that will also be used in the next section.

We begin with a description of the sensing matrix $A$. In each experimental trial, the $m \times n$ sensing matrix $A$ is generated by randomly picking $m$ rows from the $n \times n$ discrete cosine transform (DCT) matrix. Note that the $n \times n$ DCT matrix can be created by the Matlab command `dct(eye(n))` and the norm of the matrix is always 1. Sparse signals $u$ used in our experiments are generated according to (2). In each experimental trial, a length-$n$, $s$-sparse signal (a signal having exactly $s$ nonzero components), is generated in such a way that nonzero components are given by

$$\eta_1 10^{\eta_2},$$

where $\eta_1 = \pm 1$ with probability $1/2$ and $\eta_2$ is uniformly distributed in $[0,1]$. The locations of the nonzero components are randomly permuted and are the first $s$ components of the vector by Matlab command `randperm(n)`. Clearly, the range of the magnitude of nonzero components of an $s$-sparse signal is $[0, 10^\theta]$ with the parameter $\theta$ controlling this dynamic range. An observed signal (data) is collected by $b = Au + z$, where $z$ represents a Gaussian noise.

To investigate the impact of varying the parameter $\alpha$ on the performance of Algorithm 2, we consider the configuration of $n = 2^{15}$, $m = n/2$, $s = 0.05$ and the dynamic range parameter $\theta = 5$. The observed data is noise free. Six different values of $\alpha$, namely, 0.0025, 0.005, 0.01, 0.02, 0.04, and 0.08, are tested. Figure 3.3(a) depicts the traces of the relative $\ell_1$-error (see (19)) against the number of iterations for each $\alpha$. As it can be seen from this figure, for $\alpha = 0.0025$, the smallest value in our test, the relative $\ell_1$-error drops rapidly from 1 to $10^{-4}$, stabilizes with insignificant changes for about 1200 iterations and then quickly drops again to the level of $10^{-16}$. When $\alpha$ increases from 0.0025 to 0.08, the number of iterations required for the relative $\ell_1$-error dropping from 1 to $10^{-4}$ increases. In the meanwhile, the numbers of iterations for the transitions from the first sharp jump region to the second one decrease. For example, it is about 700 for $\alpha = 0.005$ and only few iterations for $\alpha = 0.08$. These observations motivate us to extend Algorithm 2 to a scenario in which the parameter $\alpha$ can be updated during the iteration with the goal of reducing the number of iterations. The proposed approach is rather simple. It begins with a relative small $\alpha$ and then increases it for every given amount of iterations. A detailed flow of this new approach is given in Algorithm 3.

Three new parameters introduced in Algorithm 3 are integers $p > 0$, $\tau > 1$, and $T > 0$. The parameter $T$ is the allowable maximum number of updating the parameters $\alpha$ and $\beta$. For
Figure 3.3: The relative $\ell_1$ error (the vertical axis) versus number of iterations (the horizontal axis). (a) Convergence of Algorithm 2 with different values of $\alpha$; (b) Convergence of Algorithm 3.

**Algorithm 3** (A variant of Algorithm 1 for model (BP$_\epsilon$))

Given: integers $p > 0$, $\tau > 1$, and $T > 0$; $\epsilon > 0$

Initialization: $v^0 \in \mathbb{R}^m$, $u^0 \in \mathbb{R}^n$, $\alpha > 0$, and $\beta > 0$ with $\frac{\beta}{\alpha} < \frac{1}{\|A\|}$; set $v^{-1} = v^0 - (Au^0 - d^0)$

repeat ($k \geq 0$)

Step 1: Compute $u^{k+1}$ using Step 1 of Algorithm 2

Step 2: Compute $v^{k+1}$ using Step 2 of Algorithm 2

Step 3: If $k$ is a multiple of $p$ and the number of changing the parameters $\alpha$ and $\beta$ does not exceed $T$, update

$$
\alpha \leftarrow \tau \alpha, \quad \beta \leftarrow \frac{\beta}{\tau}
$$

until a given stopping criteria is met
each update, the pair \((\alpha, \beta)\) will change to \((\tau \alpha, \beta/\tau)\) that will keep the ratio \(\frac{\beta}{\alpha}\) unchanged. The parameter \(p\) is to indicate that the underlying algorithm with a pair \((\alpha, \beta)\) will iterate \(p\) times before the algorithm with the pair \((\tau \alpha, \beta/\tau)\) runs another \(p\) times. We now demonstrate the efficiency of varying the parameters \(\alpha\) and \(\beta\) via applying Algorithm 3 for the same data used in Figure 3.3(a). We set \(T = 6\), \(\tau = 4\), and \(p = 20\) and initialize \(\alpha = \frac{10^2}{n \|A^\top b\|_\infty}\). Again, we choose \(\beta\) by using (17). The corresponding result is shown in Figure 3.3(b). It is clear to see that it takes about 200 iterations to drop the relative \(\ell_1\) error down below \(10^{-14}\). Hence the strategy of updating the parameters \(\alpha\) and \(\beta\) as described in Algorithm 3 is reasonable. Extensive experiments using this algorithm for problems \((BP)\) and \((BP_\epsilon)\) will be presented in the next section.

4 Numerical Simulations

This section is devoted to showing the numerical performance of the proposed Algorithm 2 for compressive sampling. We use the algorithm NESTA [2] as a comparison. We focus on sparse signals with various dynamic ranges and measurement matrices from randomly partial discrete cosine transforms (DCTs) and evaluate performance of algorithms in terms of various error metrics, speed, and robustness to noise. All the experiments are performed in Matlab 7.11 on Thinkpad T400 with Intel Core Duo CPU @2.26G, 3GB RAM on Windows Vista Home Basic operating system.

In our experiments, the \(m \times n\) sensing matrix \(A\) and length-\(n\) and \(s\)-sparse signals are generated by the same way as we did in the previous section. That is, the matrix \(A\) is generated by randomly picking \(m\) rows from the \(n \times n\) discrete cosine transform (DCT) matrix while the length-\(n\) and \(s\)-sparse signals are created by (18).

The accuracy of a solution obtained from a specific algorithm is quantified by the relative \(\ell_2\)-error, the relative \(\ell_1\)-error, and the absolute \(\ell_\infty\)-error defined, respectively, as follows:

\[
\frac{\|u - u_\circ\|}{\|u\|}, \quad \frac{||u||_1 - ||u_\circ||_1}{||u||_1}, \quad ||u - u_\circ||_\infty, \tag{19}
\]

where \(u\) is the true data and \(u_\circ\) is the restored data. All results reported in the tables of this section are the means and the standard deviations of these relative errors from simulations that were performed 20 trials.

The rest of this section consists of two parts. Part one contains the results for \((BP)\) while Part two contains the results for \((BP_\epsilon)\).

In Part one, we compare the performance of Algorithm 3 with that of the NESTA [2]. The algorithm NESTA was developed by applying a smoothing technique for the nonsmooth \(\ell_1\)-norm and an accelerated first-order scheme, both from Nesterov’s work [21]. A parameter denoted by \(\mu\) is used to control how close the smoothed \(\ell_1\)-norm to the \(\ell_1\)-norm will be. Two different levels of \(\mu\), namely, \(\mu = 10^{-2}\) and \(\mu = 10^{-7}\), are used for the NESTA in experiments. A parameter \(\delta\) for tolerance in the NESTA varies for different values of the smoothing parameter \(\mu\) and needs to be determined. We finally choose \(\delta = 10^{-8}\) for \(\mu = 10^{-2}\) and \(\delta = 10^{-12}\) for \(\mu = 10^{-7}\) since such choices lead to reasonable results. For Algorithm 3, we set \(p = 20\) and \(T\) to be the smallest integer that is greater than \(\log_{10}(\frac{\mu}{\delta} \|A^\top b\|_\infty)\). In our experiments, we notice that \(\frac{\mu}{\delta} \|A^\top b\|_\infty\) is approximately equal to \(10^{\theta}\). As a result, \(T\) is about \(\theta + 1\). The stopping criterion of Algorithm 3 is that the relative errors between the successive iterates of the reconstructed signal should satisfy the inequality \(\|u^{k+1} - u^k\|/\|u^k\| < 10^{-15}\).

In our experiments for problem \((BP)\), the dimensions \(n\) and the dynamic ranges \(\theta\) of the unknown signals are chosen from \(\{2^{13}, 2^{15}, 2^{17}\}\) and \(\{1, 3, 5\}\), respectively. The number of nonzero entries \(s\)
is set to be 0.05n, 0.02n, 0.01n respectively for the number of measurements $m = n/2$, $n/4$, $n/8$. Figure 4.4 shows the results of Algorithm 3 and the NESTA when the number of measurements $m$ is $n/2$. The dimensions $n$ of the tested signals (from left column to right column) are $2^{13}$, $2^{15}$, $2^{17}$, respectively. Solid line with ‘o’, dashed line with ‘x’, and dashdot line with ‘+’ denote the results produced by Algorithm 3, the NESTA with $\mu = 2^{-2}$, and the NESTA with $\mu = 2^{-7}$, respectively. Three error metrics, namely $\ell_2$-error, $\ell_1$-error, $\ell_\infty$-error, are respectively displayed in the first three rows with a base 10 logarithmic scale plot for the vertical axis. We can clearly see that Algorithm 3 is significantly better than the NESTA in terms of these error metrics for the unknown signals with various dimensions and dynamic ranges. The CPU time consumed by the algorithms is reported in the last row of Figure 4.4 and are plotted by using Matlab command \texttt{errorbar}. For clarity, we shift the horizontal axis to left by 0.15 for the NESTA with $\mu = 10^{-2}$ and to right by 0.15 the NESTA with $\mu = 10^{-7}$. We can see that the CPU time for Algorithm 3 is about half of that for the NESTA with $\mu = 10^{-7}$. We also observe that the $\ell_2$-error and $\ell_1$-error of the results recovered by Algorithm 3 along with the CPU time consumed are quite robust with respect to the dimensions and the dynamic ranges of the unknown signals. The results with $m = n/4$ and $m = n/8$ are reported in Figure 4.5 and Figure 4.6, respectively. The same conclusions drawn on the case of $m = n/2$ apply to the cases of $m = n/4$ and $m = n/8$ as well.

Part two of this section is about the results for problem (BP$_r$). The settings of dimension, sparsity, and dynamic range of unknown signals for problem (BP$_r$) are the same as those for problem (BP). The only difference is that measurements in Part two are contaminated by noise. In our experiments, noise levels in the measurements vary with the dynamic ranges of the unknown signals. More precisely, the noise levels $\sigma$ are set to be 0.05, 1.0, and 5.0 corresponding to the dynamic range parameter $\theta$ of 1, 3, and 5, respectively. It turns out that the noise power is $\sigma^2 = m\sigma^2$. The stopping criteria for our algorithm is that the relative errors between the successive iterates of the reconstructed signal should satisfy the inequality $\|u^{k+1} - u^k\|/\|u^k\| < 10^{-5}$. For the smoothing parameter $\mu$ in the NESTA, we choose the default setting $\mu = \max\{0.1\sigma, 0.01\}$. Instead of presenting experimental results using figures as we did for problem (BP) in Part one, we report the results in Tables I, II, III for $m = n/2$, $m = n/4$, and $m = n/8$, respectively. From these tables, we can see that the performance of our algorithm is comparable to that of the NESTA in terms of “relative $\ell_2$-error”. There is up to 1 order of improvement of accuracy regarding the relative $\ell_1$-error” with Algorithm 3. The $\ell_\infty$-error by Algorithm 3 is always lower than that by the NESTA. Regarding the consumed CPU time, Algorithm 3 consumes at most half of the CPU time that the NESTA does.

5 Conclusions

In this paper, we have reformulated $\ell_1$-norm minimization problems (BP) and (BP$_r$) via indicator functions as unconstrained minimization problems. The objective function for each unconstrained problem is the sum of the $\ell_1$-norm of the underlying signal $u$ and the indicator function of a set in $\mathbb{R}^m$, which is $\{0\}$ for (BP) or the $\epsilon$-ball for (BP$_r$), composing with the affine transformation $Au - b$. Due to the structure of this objective function and the availability of the explicit forms of the proximity operators for both the $\ell_1$-norm and the indicator function, an accurate and efficient algorithm is developed for recovering sparse signals. The algorithm outperforms the state-of-the-art algorithm NESTA in terms of the relative $\ell_2$, the relative $\ell_1$, and the absolute $\ell_\infty$ error measures as well as the CPU time for tested signals ranging from a low dynamic range to a high dynamic range with different sizes.
Figure 4.4: The $\ell_2$-error (row 1), $\ell_1$-error (row 2), $\ell_\infty$-error (row 3), and the CPU time (row 4) as functions of magnitude of the dynamic range $[0, 10^\theta]$, where $\theta$ is the $x$-axis parameter. Solid line: Algorithm 3. Dashed line: NESTA with $\mu = 2^{-2}$. Dashdot line: NESTA with $\mu = 2^{-7}$. The dimensions $n$ of the tested signals (from left column to right column) are $2^{13}$, $2^{15}$, $2^{17}$, respectively. The corresponding numbers of nonzero samples in the tested signals (from left column to right column) are $0.05 \times 2^{13}$, $0.02 \times 2^{15}$, $0.01 \times 2^{17}$, respectively. The number of measurements $m$ is $n/2$. 
Figure 4.5: The $\ell_2$-error (row 1), $\ell_1$-error (row 2), $\ell_\infty$-error (row 3), and the CPU time (row 4) as functions of magnitude of the dynamic range $[0, 10^\theta]$, where $\theta$ is the $x$-axis parameter. Solid line: Algorithm 3. Dashed line: NESTA with $\mu = 2^{-2}$. Dashdot line: NESTA with $\mu = 2^{-7}$. The dimensions $n$ of the tested signals (from left column to right column) are $2^{13}$, $2^{15}$, $2^{17}$, respectively. The corresponding numbers of nonzero samples in the tested signals (from left column to right column) are $0.05 \times 2^{13}$, $0.02 \times 2^{15}$, $0.01 \times 2^{17}$, respectively. The number of measurements $m$ is $n/4$. 
Figure 4.6: The $\ell_2$-error (row 1), $\ell_1$-error (row 2), $\ell_\infty$-error (row 3), and the CPU time (row 4) as functions of magnitude of the dynamic range $[0, 10^\theta]$, where $\theta$ is the $x$-axis parameter. Solid line: Algorithm 3. Dashed line: NESTA with $\mu = 2^{-2}$. Dashdot line: NESTA with $\mu = 2^{-7}$. The dimensions $n$ of the tested signals (from left column to right column) are $2^{13}, 2^{15}, 2^{17}$, respectively. The corresponding numbers of nonzero samples in the tested signals (from left column to right column) are $0.05 \times 2^{13}, 0.02 \times 2^{15}, 0.01 \times 2^{17}$, respectively. The number of measurements $m$ is $n/8$. 
<table>
<thead>
<tr>
<th>Method</th>
<th>relative $\ell_2$-error</th>
<th>relative $\ell_1$-error</th>
<th>absolute $\ell_\infty$-error</th>
<th>CPU time(s)</th>
</tr>
</thead>
<tbody>
<tr>
<td>proposed</td>
<td>(3.60e-2, 1.42e-3)</td>
<td>(3.92e-3, 1.22e-3)</td>
<td>(3.66e-1, 3.56e-2)</td>
<td>(3.49e-1, 1.35e-2)</td>
</tr>
<tr>
<td>NESTA</td>
<td>(4.19e-2, 1.28e-3)</td>
<td>(1.96e-2, 8.60e-4)</td>
<td>(4.23e-1, 3.64e-2)</td>
<td>(1.47, 9.27e-2)</td>
</tr>
<tr>
<td>n = 2</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>proposed</td>
<td>(3.59e-2, 7.48e-4)</td>
<td>(3.59e-3, 8.20e-4)</td>
<td>(3.93e-1, 2.64e-2)</td>
<td>(1.49, 4.22e-2)</td>
</tr>
<tr>
<td>NESTA</td>
<td>(4.20e-2, 7.52e-4)</td>
<td>(1.88e-2, 7.23e-4)</td>
<td>(4.54e-1, 3.04e-2)</td>
<td>(4.83, 1.90e-1)</td>
</tr>
<tr>
<td>n = 2</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>proposed</td>
<td>(3.58e-2, 3.24e-4)</td>
<td>(3.44e-3, 4.48e-4)</td>
<td>(4.16e-1, 1.92e-2)</td>
<td>(8.96, 1.00e-1)</td>
</tr>
<tr>
<td>NESTA</td>
<td>(4.18e-2, 3.79e-4)</td>
<td>(1.84e-2, 3.21e-4)</td>
<td>(4.89e-1, 2.75e-2)</td>
<td>(2.72e+1, 3.07e-1)</td>
</tr>
<tr>
<td>n = 2</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>proposed</td>
<td>(1.18e-2, 8.60e-4)</td>
<td>(1.14e-3, 6.52e-4)</td>
<td>(7.13, 5.06e-1)</td>
<td>(4.03e-1, 1.70e-2)</td>
</tr>
<tr>
<td>NESTA</td>
<td>(1.27e-2, 7.89e-4)</td>
<td>(1.17e-2, 1.14e-3)</td>
<td>(7.97, 4.52e-1)</td>
<td>(1.66, 2.10e-1)</td>
</tr>
<tr>
<td>n = 2</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>proposed</td>
<td>(1.16e-2, 2.32e-4)</td>
<td>(1.19e-3, 3.21e-4)</td>
<td>(7.76, 5.41e-1)</td>
<td>(1.87, 7.58e-2)</td>
</tr>
<tr>
<td>NESTA</td>
<td>(1.26e-2, 3.06e-4)</td>
<td>(1.12e-2, 4.40e-4)</td>
<td>(8.69, 8.17e-1)</td>
<td>(5.38, 1.54e-1)</td>
</tr>
<tr>
<td>n = 2</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>proposed</td>
<td>(1.15e-2, 1.24e-4)</td>
<td>(1.25e-3, 1.90e-4)</td>
<td>(8.32, 5.66e-1)</td>
<td>(1.00e+1, 8.09e-2)</td>
</tr>
<tr>
<td>NESTA</td>
<td>(1.27e-2, 1.85e-4)</td>
<td>(1.10e-2, 2.40e-4)</td>
<td>(9.10, 4.25e-1)</td>
<td>(3.14e+1, 1.03)</td>
</tr>
<tr>
<td>n = 2</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>proposed</td>
<td>(7.14e-4, 7.30e-5)</td>
<td>(5.23e-5, 4.21e-5)</td>
<td>(3.41e+1, 3.41e+1)</td>
<td>(4.69e-1, 2.70e-2)</td>
</tr>
<tr>
<td>NESTA</td>
<td>(7.83e-4, 4.74e-5)</td>
<td>(8.92e-4, 1.07e-4)</td>
<td>(3.86e+1, 2.53)</td>
<td>(1.46, 1.09e-1)</td>
</tr>
<tr>
<td>n = 2</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>proposed</td>
<td>(7.19e-4, 3.28e-5)</td>
<td>(2.59e-5, 1.81e-5)</td>
<td>(3.79e+1, 2.74)</td>
<td>(2.36, 1.13e-1)</td>
</tr>
<tr>
<td>NESTA</td>
<td>(7.95e-4, 3.47e-5)</td>
<td>(8.64e-4, 5.54e-5)</td>
<td>(4.19e+1, 2.28)</td>
<td>(4.83, 1.70e-1)</td>
</tr>
<tr>
<td>n = 2</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>proposed</td>
<td>(7.10e-4, 1.23e-5)</td>
<td>(1.33e-5, 1.06e-5)</td>
<td>(4.04e+1, 1.76)</td>
<td>(1.93e+1, 6.54e-2)</td>
</tr>
<tr>
<td>NESTA</td>
<td>(7.91e-4, 1.48e-5)</td>
<td>(8.26e-4, 2.79e-5)</td>
<td>(4.58e+1, 2.19)</td>
<td>(2.73e+1, 4.90e-1)</td>
</tr>
</tbody>
</table>

Table I: The columns “relative $\ell_2$-error”, “relative $\ell_1$-error”, and “absolute $\ell_\infty$-error” display the errors between recovered signals and the true data. The column “CPU time” shows CPU time of algorithms; The number of measurements $m$ is $n/2$ and the test signals are $s$-sparse with $s = 0.05n$. The pair ($\cdot$, $\cdot$) represents the mean and the standard deviation.

References


<table>
<thead>
<tr>
<th>Method</th>
<th>relative $\ell_2$-error</th>
<th>relative $\ell_1$-error</th>
<th>absolute $\ell_\infty$-error</th>
<th>CPU time(s)</th>
</tr>
</thead>
<tbody>
<tr>
<td>proposed</td>
<td>$n = 2^{13}$</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td>$(6.03e-2, 3.27e-3)$</td>
<td>$(5.88e-3, 3.65e-4)$</td>
<td>$(5.36e-1, 5.39e-2)$</td>
<td>$(3.85e-1, 4.70e-2)$</td>
</tr>
<tr>
<td>NESTA</td>
<td>$(7.27e-2, 2.96e-3)$</td>
<td>$(6.77e-1, 6.80e-2)$</td>
<td>$(1.69, 9.79e-2)$</td>
<td></td>
</tr>
<tr>
<td></td>
<td>$n = 2^{15}$</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td>$(6.12e-2, 2.13e-3)$</td>
<td>$(6.38e-3, 1.72e-3)$</td>
<td>$(6.20e-1, 4.35e-2)$</td>
<td>$(1.67, 6.27e-2)$</td>
</tr>
<tr>
<td>NESTA</td>
<td>$(7.18e-2, 2.07e-3)$</td>
<td>$(7.17e-1, 3.09e-2)$</td>
<td>$(6.54, 1.82e-1)$</td>
<td></td>
</tr>
<tr>
<td></td>
<td>$n = 2^{17}$</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td>$(6.11e-2, 1.98e-3)$</td>
<td>$(6.73e-3, 8.89e-4)$</td>
<td>$(6.55e-1, 4.91e-2)$</td>
<td>$(9.69, 1.54e-1)$</td>
</tr>
<tr>
<td>NESTA</td>
<td>$(7.17e-2, 1.25e-3)$</td>
<td>$(7.52e-1, 3.65e-2)$</td>
<td>$(3.28e+1, 3.03e-1)$</td>
<td></td>
</tr>
<tr>
<td></td>
<td>$a = 1, \sigma = 0.05$</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td>$(1.89e-2, 2.11e-3)$</td>
<td>$(1.48e-3, 8.76e-4)$</td>
<td>$(1.01e+1, 9.45e-1)$</td>
<td>$(4.76e-1, 6.04e-2)$</td>
</tr>
<tr>
<td>NESTA</td>
<td>$(2.02e-2, 1.85e-3)$</td>
<td>$(1.20e+1, 7.92e-1)$</td>
<td>$(1.90, 1.10e-1)$</td>
<td></td>
</tr>
<tr>
<td></td>
<td>$n = 2^{15}$</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td>$(1.68e-2, 1.10e-3)$</td>
<td>$(1.43e-3, 5.52e-4)$</td>
<td>$(1.15e+1, 9.40e-1)$</td>
<td>$(2.46, 7.08e-2)$</td>
</tr>
<tr>
<td>NESTA</td>
<td>$(2.06e-2, 9.19e-4)$</td>
<td>$(1.30e+1, 7.85e-1)$</td>
<td>$(7.39, 1.41e-1)$</td>
<td></td>
</tr>
<tr>
<td></td>
<td>$n = 2^{17}$</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td>$(1.88e-2, 4.86e-4)$</td>
<td>$(1.56e-3, 4.11e-4)$</td>
<td>$(1.24e+1, 8.19e-1)$</td>
<td>$(1.41e+1, 1.68e-1)$</td>
</tr>
<tr>
<td>NESTA</td>
<td>$(2.04e-2, 5.61e-4)$</td>
<td>$(1.36e+1, 6.71e-1)$</td>
<td>$(3.67e+1, 4.38e-1)$</td>
<td></td>
</tr>
<tr>
<td></td>
<td>$a = 3, \sigma = 1$</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td>$(1.69e-3, 3.61e-4)$</td>
<td>$(8.85e-5, 6.74e-5)$</td>
<td>$(5.01e+1, 5.52)$</td>
<td>$(4.89e+1, 2.33e-2)$</td>
</tr>
<tr>
<td>NESTA</td>
<td>$(1.25e-3, 1.29e-4)$</td>
<td>$(5.92e-1, 5.25)$</td>
<td>$(1.84, 1.22e-1)$</td>
<td></td>
</tr>
<tr>
<td></td>
<td>$n = 2^{15}$</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td>$(1.19e-3, 7.33e-5)$</td>
<td>$(6.92e-5, 4.46e-5)$</td>
<td>$(6.49e+1, 3.92)$</td>
<td>$(2.28, 1.36e-1)$</td>
</tr>
<tr>
<td>NESTA</td>
<td>$(1.27e-3, 8.32e-5)$</td>
<td>$(6.32e+1, 3.96)$</td>
<td>$(7.18, 1.55e-1)$</td>
<td></td>
</tr>
<tr>
<td></td>
<td>$n = 2^{17}$</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td>$(1.15e-3, 3.86e-5)$</td>
<td>$(5.29e-5, 3.72e-5)$</td>
<td>$(6.11e+1, 5.45)$</td>
<td>$(1.33e+1, 6.76e-1)$</td>
</tr>
<tr>
<td>NESTA</td>
<td>$(1.29e-3, 4.24e-5)$</td>
<td>$(6.94e+1, 4.63)$</td>
<td>$(3.48e+1, 1.03)$</td>
<td></td>
</tr>
</tbody>
</table>

Table II: The columns “relative $\ell_2$-error”, “relative $\ell_1$-error”, and “absolute $\ell_\infty$-error” display the errors between recovered signals and the true data. The column “CPU time” shows CPU time of algorithms; The number of measurements $m$ is $n/4$ and the test signals are $s$-sparse with $s = 0.02n$. The pair $(\cdot, \cdot)$ represents the mean and the standard deviation.
<table>
<thead>
<tr>
<th>Method</th>
<th>$n = 2^{13}$</th>
<th>$n = 2^{15}$</th>
<th>$n = 2^{17}$</th>
</tr>
</thead>
<tbody>
<tr>
<td>proposed</td>
<td>(1.99e-1, 9.92e-3)</td>
<td>(1.97e-1, 1.00e-3)</td>
<td>(1.94e-1, 2.39e-3)</td>
</tr>
<tr>
<td></td>
<td>(2.00e-2, 6.99e-3)</td>
<td>(2.00e-2, 7.00e-3)</td>
<td>(2.00e-2, 8.96e-4)</td>
</tr>
<tr>
<td></td>
<td>(5.67e-1, 7.56e-2)</td>
<td>(5.67e-1, 7.56e-2)</td>
<td>(5.67e-1, 7.56e-2)</td>
</tr>
<tr>
<td>NESTA</td>
<td>(1.22e-1, 1.16e-2)</td>
<td>(1.22e-1, 1.16e-2)</td>
<td>(1.22e-1, 1.16e-2)</td>
</tr>
<tr>
<td></td>
<td>(3.18e-2, 6.03e-3)</td>
<td>(3.18e-2, 6.03e-3)</td>
<td>(3.18e-2, 6.03e-3)</td>
</tr>
<tr>
<td></td>
<td>(9.88e-1, 9.74e-2)</td>
<td>(9.88e-1, 9.74e-2)</td>
<td>(9.88e-1, 9.74e-2)</td>
</tr>
<tr>
<td></td>
<td>(1.98, 1.55e-1)</td>
<td>(1.98, 1.55e-1)</td>
<td>(1.98, 1.55e-1)</td>
</tr>
<tr>
<td>proposed</td>
<td>(1.03e-1, 5.06e-3)</td>
<td>(1.03e-1, 5.06e-3)</td>
<td>(1.03e-1, 5.06e-3)</td>
</tr>
<tr>
<td></td>
<td>(1.00e-2, 3.46e-3)</td>
<td>(1.00e-2, 3.46e-3)</td>
<td>(1.00e-2, 3.46e-3)</td>
</tr>
<tr>
<td></td>
<td>(9.46e-1, 7.05e-2)</td>
<td>(9.46e-1, 7.05e-2)</td>
<td>(9.46e-1, 7.05e-2)</td>
</tr>
<tr>
<td></td>
<td>(8.39, 2.17e-1)</td>
<td>(8.39, 2.17e-1)</td>
<td>(8.39, 2.17e-1)</td>
</tr>
<tr>
<td>NESTA</td>
<td>(1.22e-1, 3.83e-3)</td>
<td>(1.22e-1, 3.83e-3)</td>
<td>(1.22e-1, 3.83e-3)</td>
</tr>
<tr>
<td></td>
<td>(2.66e-2, 2.18e-3)</td>
<td>(2.66e-2, 2.18e-3)</td>
<td>(2.66e-2, 2.18e-3)</td>
</tr>
<tr>
<td></td>
<td>(1.09, 8.11e-2)</td>
<td>(1.09, 8.11e-2)</td>
<td>(1.09, 8.11e-2)</td>
</tr>
<tr>
<td></td>
<td>(8.39, 2.17e-1)</td>
<td>(8.39, 2.17e-1)</td>
<td>(8.39, 2.17e-1)</td>
</tr>
</tbody>
</table>

Table III: The columns “relative $\ell_2$-error”, “relative $\ell_1$-error”, and “absolute $\ell_\infty$-error” display the errors between recovered signals and the true data. The column “CPU time” shows CPU time of algorithms; The number of measurements $m$ is $n/8$ and the test signals are $s$-sparse with $s = 0.01n$. The pair $(\cdot, \cdot)$ represents the mean and the standard deviation.
