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The Alternate-Block-Factorization Procedure for Systems of Partial Differential Equations

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Abstract

The alternate-block-factorization (ABF) method is a procedure for partially decoupling systems of elliptic partial differential equations by means of a carefully chosen change of variables. By decoupling, we mean that the ABF strategy attempts to reduce intra-equation coupling in the system rather than intra-grid coupling for a single elliptic equation in the system. This has the effect of speeding convergence of commonly used iteration schemes, which use the solution of a sequence of linear elliptic PDEs as their main computational step. Algebraically, the change of variables is equivalent to a postconditioning of the original system. The results of using ABF postconditioning on some problems arising from semiconductor device simulation are discussed.

1 Introduction

In this paper, we are interested in approximately solving a system of elliptic partial differential equations (PDEs) on a domain $\Omega \in \mathbb{R}^d$ with appropriate boundary conditions. Let us write the system in terms of scalar PDEs as follows

$$(1) \quad L(z) = 0$$

$$\begin{pmatrix} L_1(z_1, z_2, \dots, z_m) \\ L_2(z_1, z_2, \dots, z_m) \\ \vdots \\ L_m(z_1, z_2, \dots, z_m) \end{pmatrix}$$

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Equations (3)-(5), or (8)-(10), ignore complications like variable mobilities, oxide regions, and the generation and recombination of carriers. However,

$$L_3(u, v, w) = -\Delta \cdot \Delta (e^{u-v} w) = 0, \tag{10}$$

$$L_2(u, v, w) = \Delta \cdot \Delta (e^{u-v} v) = 0, \tag{9}$$

$$L_1(u, v, w) = -\Delta^2 u + e^{u-v} - e^{v-u} - N(x) = 0, \tag{8}$$

In these variables, the equations become

$$\begin{aligned} n &= e^{u-v}, \\ d &= e^{v-u}. \end{aligned} \tag{7}$$

Equations (3)-(5) are sometimes written in terms of *quasi-Fermi* variables, u, v , and w , defined by

There are one carrier variants of this problem that amount to dropping (5) and p or, sometimes, (4) and n . We assume that Ω is a simply connected polygonal region. $N(x)$ is the given doping profile. The unknown dependent variables are the electrostatic potential, u , and electron and hole densities, n and p , respectively; we refer to u, n , and p as the *primitive* variables. Dirichlet boundaries conditions are normally given on part of the boundary, $\partial\Omega_1 \subset \partial\Omega$, with homogeneous Neumann conditions elsewhere.

$$L_3(u, n, p) = -\Delta \cdot \Delta (p \Delta u + \Delta p) = 0, \tag{5}$$

$$L_2(u, n, p) = \Delta \cdot \Delta (n \Delta u - \Delta n) = 0, \tag{4}$$

$$L_1(u, n, p) = -\Delta^2 u + n - p - N(x) = 0, \tag{3}$$

Our particular interest has been in the system of $m = 3$ drift-diffusion equations that occur in semiconductor device modeling [6, 8, 23]. A simplified form of this system is given by

If there are ν degrees of freedom associated with the underlying discrete approximation to each PDE, then (2) represents $m\nu$ nonlinear algebraic equations to be solved.

$$L_\nu(z) = \begin{pmatrix} L_1(z_1^1, z_2^1, \dots, z_\nu^1) \\ \vdots \\ L_m(z_1^m, z_2^m, \dots, z_\nu^m) \end{pmatrix} = 0. \tag{2}$$

where $z(x) \in \mathbb{R}^m$ is a vector function. Systems of this type arise from many applications in science and engineering. In general, we are interested in a discretized form of (1). Let Ω^h be a triangulation of Ω . Then a finite-difference, finite-element, or finite-volume method can be applied to (1) resulting in

Both the plug-in and coupled methods are widely used in semiconductor device simulators (for example, see [6, 23, 21, 4, 13, 3]). The plug-in method results in smaller, mathematically more tractable systems of equations. If the PDES in (2) are weakly coupled, the convergence can be quite rapid. When the PDES are strongly coupled, the convergence of the outer plug-in iteration can be quite slow, or even diverge. For the semiconductor problem in quasi-Fermi

$$L' = \left(\frac{\partial L}{\partial z} \right) = \begin{bmatrix} L_{11} & L_{12} & \dots & L_{1m} \\ L_{21} & L_{22} & \dots & L_{2m} \\ \vdots & \vdots & \ddots & \vdots \\ L_{m1} & L_{m2} & \dots & L_{mm} \end{bmatrix} \quad (14)$$

where L' is a (possibly approximate) Jacobian and d^k is a suitably chosen scalar damping factor [5]. L' can be viewed as a block matrix

$$L'(z^k)x^k = -L(z^k), \quad (12)$$

$$z^{k+1} = z^k + d^k x^k \quad (13)$$

The second approach is to linearize (2) and apply Newton's method. Each step of this method requires solving a system of linear equations and then updating

the Gauss-Seidel-Newton method. Each method, as we will shortly discuss; we refer to the resulting overall scheme as the plug-in method. One common scheme for solving each of the m scalar PDES is Newton's

$$L_i(z_{k+1}^1, z_k^2, \dots, z_{k+1}^i, z_k^{i+1}, z_k^{i+2}, \dots, z_k^m) = 0 \quad (11)$$

A variety of approaches exists for dealing with the nonlinear equations represented by (2), but we will concentrate on the two principle techniques used in device simulators. For notational simplicity, we will drop the use of the superscript i to denote the discrete forms since usually only the discretized systems will be considered in what follows. The first method for solving (2) is nonlinear Gauss-Seidel [20], known as Gummel's iteration in the device-simulation literature [9]. Suppose an initial guess $(z^0)^T = (z_0^1, z_0^2, \dots, z_0^m)$ is given. To go from the k th iterate, z^k , to the $(k+1)$ st iterate, z^{k+1} , simply solve, for $i = 1, 2, \dots, m$, the i th PDE

usually employed. Methods taking advantage of the convective-diffusive nature of the PDES are usually employed. Similar discretizations can be applied to (4) and (5), but specialized methods taking advantage of the convective-diffusive nature of the PDES are usually employed. In our main interest — solving discrete analogs of (3)-(5). Equation (3) can be discretized by a number of finite-difference, finite-element, or finite-volume methods. (3)-(5) capture some of the difficulties that occur in practice and let us focus

variables, (8-10), the plug-in method converges well when the iterates are far from the solution [14]; device simulation programs often use a few initial plug-in iterations to improve an initial guess before switching to a coupled approach. The coupled method takes the interactions between the equations into account (though the off-diagonal blocks of L'). However, dealing with (12) can be arduous because the entire linearized system must be computed, assembled, and solved.

Equation (12) can itself be solved by an inner iteration. An obvious approach is to employ a block Gauss-Seidel iteration, which would require solving systems involving the diagonal blocks, L_{ii} , of the Jacobian; we refer to this approach as the Newton-Gauss-Seidel method. These linear equations would be the same as those solved in the plug-in method if Newton's method were used to deal with the scalar PDEs represented by (11). Hence, in the case of tightly coupled PDEs, the slow convergence of the outer iteration of the plug-in approach will correspond to slow convergence of the inner block iteration for the coupled method.

In both the plug-in and coupled Newton-Gauss-Seidel approaches, the idea is to reduce the solution of a large coupled system into a series of discrete scalar elliptic PDEs. Such PDEs are fairly well understood and a wide variety of algorithms exists to solve them. (Some of the possible algorithms include sparse direct methods, Krylov-subspace iterations, multigrid techniques, and fast Poisson solvers.)

The *alternate-block-factorization* (ABF) technique, which we are about to discuss in detail, can be applied in the context of either the plug-in or coupled iterations, assuming the coupled method makes use of an inner block iteration. (The ABF idea and some results were partially described in [3].) The ABF method was motivated by trying to find a 'better' sequence of scalar problems to solve in order to speed convergence of an outer iteration. This is done by a temporary local change of variables. Algebraically, this amounts to preconditioning (12) on the right (postconditioning it).

The ABF strategy is attractive because it is simple to describe and implement as well as having limited storage requirements. The computational work required, mainly the inversion of ν matrices of order m and the diagonal scaling of matrices corresponding to scalar problems, is generally a lower-order term in the overall work estimate. The ABF approach can substantially reduce the total effort to solve a tightly coupled system.

The remainder of this paper is organized as follows. In § 2, the ABF technique will be derived. Some special cases will be studied in § 3. The results of some computational experiments will be shown in § 4. Some conclusions will be drawn in § 5.

2 The ABF Method

The alternate-block-factorization (ABF) method can be motivated intuitively through the following formal line of reasoning. The plug-in iteration (11) converges slowly or diverges when z_k^i is strongly coupled to the other unknown functions through their PDEs, and thus not well determined by its own PDE. This suggests that we have somehow made a 'bad' choice in assigning unknowns to equations; it is quite possible that none of the unknown functions is a 'good' choice. This leads us to (temporarily) associate a new unknown function $\zeta_i = (\zeta_i^1, \zeta_i^2, \dots, \zeta_i^m)$ with the i th PDE. Our goal is to choose the new unknowns such that they can be well determined by the solution of their associated PDEs; we do this by attempting to weaken the coupling between the PDEs. Formally linearizing (1) about the current iterate z_k^i , we have, for $i = 1, 2, \dots, m$,

$$(15) \quad \sum_{j=1}^m \left(\frac{\partial L_i}{\partial z_j^i} \right) (z_k^i) \left\{ \frac{\partial \zeta_j^i}{\partial \zeta_i^i} \delta \zeta_i^i = -L_i(z_k^i) \right\}.$$

When discretized, each $\partial L_i / \partial z_j^i$ becomes a $\nu \times \nu$ matrix, the i j block of the Jacobian, L_{ij} . The terms $\partial z_j^i / \partial \zeta_i^i$ generally must be considered dense $\nu \times \nu$ matrices. However, if we assume that the coupling between the PDEs is largely localized by grid point, then it is reasonable to approximate all the $\partial z_j^i / \partial \zeta_i^i$ terms by diagonal matrices; these matrices are estimated by solving local problems at each grid point, holding the solution at all other grid points fixed. We will now show how this procedure can be performed in practice using local matrix operations.

A description of the ABF approach in terms of matrices is rather straightforward. To simplify things, let us discuss the ABF method assuming that (2) is being solved by the coupled (Newton) approach with a block Gauss-Seidel iteration being used for (12). Hence, we are interested in solving a linear system of algebraic equations

$$(16) \quad Ax = b \in \mathbb{R}^{m\nu}.$$

The matrix, A , can be written in block form as

$$(17) \quad A = \begin{bmatrix} A_{11} & A_{12} & \dots & A_{1m} \\ A_{21} & A_{22} & \dots & A_{2m} \\ \vdots & \vdots & \ddots & \vdots \\ A_{m1} & A_{m2} & \dots & A_{mm} \end{bmatrix}$$

and $A_{ij} \in \mathbb{R}^{\nu \times \nu}$. Let

$$(18) \quad D = \begin{bmatrix} D_{11} & D_{12} & \dots & D_{1m} \\ D_{21} & D_{22} & \dots & D_{2m} \\ \vdots & \vdots & \ddots & \vdots \\ D_{m1} & D_{m2} & \dots & D_{mm} \end{bmatrix}$$

Obviously, $D^{-1}v$ can be computed using dense matrix techniques since it consists of $m \times m$ matrices on its diagonal (and we are assuming that $m \ll \nu$). Thus, the ABR-postconditioned matrix, AD^{-1} , can be formed by local operations; by construction, D^{-1} eliminates the diagonals of the off-diagonal blocks of A as we will now see with an example.

$$D = PDP^T. \quad (22)$$

Once again, note that A is the matrix blocked by PDE while \bar{A} is alternately blocked by grid point. Let

$$\bar{A}_{ij} = \begin{bmatrix} (A_{11})_{ij} & (A_{12})_{ij} & \dots & (A_{1m})_{ij} \\ (A_{21})_{ij} & (A_{22})_{ij} & \dots & (A_{2m})_{ij} \\ \vdots & \vdots & \ddots & \vdots \\ (A_{m1})_{ij} & (A_{m2})_{ij} & \dots & (A_{mm})_{ij} \end{bmatrix} \in \mathbb{R}^{m \times m}. \quad (21)$$

where

$$\bar{A} = PAP^T = \begin{bmatrix} \bar{A}_{11} & \bar{A}_{12} & \dots & \bar{A}_{1\nu} \\ \bar{A}_{21} & \bar{A}_{22} & \dots & \bar{A}_{2\nu} \\ \vdots & \vdots & \ddots & \vdots \\ \bar{A}_{\nu 1} & \bar{A}_{\nu 2} & \dots & \bar{A}_{\nu\nu} \end{bmatrix} \quad (20)$$

P , such that

Let us now consider how A and \bar{A} are related. There is a permutation matrix, P , matrices reflect the coupling between the PDEs. block matrix now represents the connectivity of the mesh while the individual points in the mesh (or degrees of freedom). In this case, the reordered matrix, \bar{A} , There is an *alternate blocking* that associates matrix blocks with individual the block matrix, A , represents the coupling between the PDEs. the connectivity of the underlying spatial mesh, Ω^h . The nonzero structure of PDE blocking, the nonzero structure of the individual matrices, A_{ij} , captures are associated with PDEs, that is, A_{ij} corresponds to the i th PDE. With this resented by (16). The first is the one given by (17) where the matrix blocks

There are two natural blockings (ordering of variables) for the system representation [3, 2]. iteration can be solved via sparse direct or preconditioned iterative methods D^{-1} is the ABR postconditioner. We then solve (19) by block Gauss-Seidel (or block SSOR) iteration; the diagonal block systems arising in the Gauss-Seidel

$$(AD^{-1})(Dx) = b. \quad (19)$$

where $D_{ij} = \text{diag}(A_{ij})$. If D^{-1} exists, we postcondition (16) by

Our goal is to solve (16) via a postconditioned block Gauss-Seidel iteration; with this in mind, let us carefully examine AD^{-1} in (28). We first consider the diagonal blocks, which correspond to scalar linear elliptic PDEs. In terms of

$$\delta = (-\text{diag}(\Delta)\text{diag}(C) + \text{diag}(M))^{-1}, \quad (29)$$

with

$$AD^{-1} = \begin{bmatrix} (-\Delta\text{diag}(C) + \text{diag}(M))\delta & (-M\text{diag}(C) + C\text{diag}(M))\delta \\ (-\Delta\text{diag}(\Delta) + \Delta)\delta & (-C\text{diag}(\Delta) + M)\delta \end{bmatrix}, \quad (28)$$

was used to simplify matters. Now AD^{-1} has the form symmetric and positive definite. Note that the identity matrix, I , in (27) could be replaced by a mass matrix if a finite-element method was used; the identity Physically, carrier densities should be positive from which it follows that M is a discretization of a convection-diffusion term so $C + C^T$ is positive definite. Here Δ is a discrete Laplacian so $-\Delta$ is symmetric and positive definite. C is

$$A = \begin{bmatrix} -\Delta & C \\ -M & I \end{bmatrix}. \quad (27)$$

matrix in (23) becomes

For the one-carrier drift-diffusion equations ((3) and (4) without p), the effect of the off-diagonal blocks. The ABF method attempts to decouple a block system by reducing the The ABF-postconditioned matrix has zeroed out the diagonals of the off-diagonal

$$\delta = (D_{11}D_{22} - D_{21}D_{12})^{-1}. \quad (26)$$

where

$$AD^{-1} = A(P^T D^{-1} P) = \begin{bmatrix} A_{11}D_{22} - A_{12}D_{21} & A_{12}D_{11} - A_{11}D_{12} \\ A_{21}D_{22} - A_{22}D_{21} & A_{22}D_{11} - A_{21}D_{12} \end{bmatrix} \begin{bmatrix} \delta & 0 \\ 0 & \delta \end{bmatrix} \quad (25)$$

Assuming D^{-1} exists, the ABF-postconditioned matrix is

$$D = \begin{bmatrix} D_{11} & D_{12} \\ D_{21} & D_{22} \end{bmatrix} = \begin{bmatrix} \text{diag}(A_{11}) & \text{diag}(A_{12}) \\ \text{diag}(A_{21}) & \text{diag}(A_{22}) \end{bmatrix} \in \mathbb{R}^{2\nu \times 2\nu}. \quad (24)$$

where each $A_{ij} \in \mathbb{R}^{\nu \times \nu}$. Let D be given by

$$\begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \end{pmatrix}. \quad (23)$$

PDEs

For specificity, consider a 2×2 block system arising from a system of two

$$(33) \quad D^{-1} = \begin{bmatrix} 2\eta h^{-2} I & & \\ -2\eta h^{-2} I & I & \\ & & (2\eta h^{-2} + \epsilon h^{-1}) I \end{bmatrix}^{-1}$$

The ABE postconditioner for this example is field, $E = -\nabla u$. This problem captures some of the character of the drift-diffusion equations written in terms of the primitive variables, that is, (3)-(5). In particular, η corresponds to a carrier density, say n , while $-\epsilon$ corresponds to the electric

$$(32) \quad S = h^{-1} [-1 \ 1 \ 0].$$

and S is the tridiagonal matrix

$$(31) \quad T = h^{-2} [-1 \ 2 \ -1],$$

Here η and ϵ are nonnegative scalars, T is the tridiagonal matrix

$$(30) \quad A = \begin{bmatrix} T & & \\ -\eta T & T + \epsilon S & \\ & & I \end{bmatrix}.$$

We consider the application of ABE to the block 2×2 matrix

down. the end of this section, we will show how the ABE postconditioner can break based on either the primitive, (3)-(5), or quasi-Fermi, (8)-(10), formulation. At equations and allow us to compare the ABE method with plug-in algorithms model problems will illuminate some of the main effects in the semiconductor and other important issues. Nevertheless, we believe that the results from these will be based on an *ad hoc* Fourier analysis that ignores boundary conditions convection-diffusion problems with equispaced meshes. Our spectral analysis ods, we will now motivate the ABE method by studying some one-dimensional [19, 22, 12, 13, 14, 15]. Since the main emphasis in this paper is on meth- been studied in great detail and with substantial mathematical rigor elsewhere alize in general. The plug-in approach for the semiconductor equations has methods. Both nonlinear and linear iterative methods can be difficult to an- Let us briefly discuss methods for analyzing pre- or post-conditioned iterative

3 Model Problem Analysis

thus speeding the convergence of the block iteration. This should have the effect of reducing the coupling between the discrete PDEs, ped, there is a fair amount of cancellation in the off-diagonal blocks of A . for solving linear systems involving these blocks converge more rapidly. As ex- more positive weight on the diagonal of $-\Delta$. This will help iterative methods ric parts of the diagonal block C has relatively less prominence and there is superior to the original diagonal blocks of (27); in particular, the nonsymmet- solving equations involving these blocks as part of a block iteration, they seem

$$Sv_k \approx \frac{h}{(1-c) + is} v_k \quad (41)$$

$$Tv_k \approx \frac{h^2}{2(1-c)} v_k, \quad (40)$$

To apply the Fourier transform, we ignore the effect of the boundary conditions and consider a discrete form of the function $\exp(it\pi x)$, where $t = \sqrt{-1}$. Setting $x_j = jh$, we define the (complex-valued) vector v_k with components $\exp(itk\pi x_j)$. Excluding the first and last rows of the matrices, we have

$$p(g) = \begin{vmatrix} ad \\ be \end{vmatrix}. \quad (39)$$

A straightforward calculation shows that the spectral radius, $p(g)$, is given by

$$g = \begin{bmatrix} a & p & e \\ a & 0 & -1 \\ 0 & 0 & b \end{bmatrix}. \quad (38)$$

and the corresponding iteration matrix

$$\begin{bmatrix} a & p \\ e & b \end{bmatrix} \quad (37)$$

First consider the 2×2 matrix

through the (formal) application of the Fourier transform [7].
because $ST \neq TS$. However, we can get an idea of the behavior of the method of (scalar) 2×2 iteration matrices. This cannot be done in a rigorous fashion. We would like to reduce this problem to that of computing the spectral radii

$$G = \begin{bmatrix} (1 + \epsilon h/2)T + \eta I & \epsilon \eta (S - hT/2) \\ 0 & (1 + \eta h^2/2)T + \epsilon S \end{bmatrix} \begin{bmatrix} -1 & 0 \\ 0 & h^2 T/2 - I \end{bmatrix}. \quad (36)$$

must consider the eigenvalues of the matrix
conditioning. To compute the spectral radius of the ABE iteration matrix, we notice that the off-diagonal blocks have zero diagonals, as a result of the post-

$$AD^{-1} = \sigma^{-1} \begin{bmatrix} (1 + \epsilon h/2)T + \eta I & \epsilon \eta (S - hT/2) \\ I - h^2 T/2 & (1 + \eta h^2/2)T + \epsilon S \end{bmatrix}. \quad (35)$$

The postconditioned matrix AD^{-1} is given by

$$\sigma = \left(\frac{h^2}{2} \right) \left(1 + \frac{\epsilon h}{2} + \frac{\eta h^2}{2} \right). \quad (34)$$

with

$$= \sigma^{-1} \begin{bmatrix} (1 + \epsilon h/2)I & \eta I \\ -h^2 I/2 & I \end{bmatrix}$$

Let us contrast these results for the ABR method with the block Gauss-Seidel iteration applied directly to A of (30). Using the Fourier transform technique

follows from (46). Finally, if all parameters except ϵ are fixed, then $\rho_{ABR} \rightarrow 0$ as $\epsilon \rightarrow 0$ or $\epsilon \rightarrow \infty$. Figure 1 shows the unimodal character of the ABR postconditioner. of ρ_{ABR} ($\partial \rho_{ABR} / \partial \eta = 0$) for some $\eta > 0$, corresponding to a case of moderate (strong coupling), convergence should be quite rapid; there is clearly a maximum from the semiconductor device model. For small η (weak coupling) or large η (or $\eta \rightarrow \infty$). This behavior with respect to η is consistent with empirical results holding all parameters except η fixed, it is clear from (46) that $\rho_{ABR} \rightarrow 0$ if $\eta \rightarrow 0$ possible to consider the behavior with respect to other parameters. For example, This, of course, is the behavior in the asymptotic limit as $h \rightarrow 0$; it is also $1 - O(h)$.

In the extreme case when $h \rightarrow 0$, $\eta \rightarrow \infty$, and $\epsilon \rightarrow \infty$, we find that $\rho_{ABR} =$

$$\rho_{ABR} \approx \left(\frac{\pi^2 + \eta}{\eta} \right) \left(\frac{\sqrt{\pi^2 + \epsilon^2}}{\epsilon} \right) > 1. \quad (49)$$

so, in the limit as $h \rightarrow 0$,

$$1 - c \approx \frac{\pi^2 h^2}{2}, \quad (47)$$

$$s \approx \pi h, \quad (48)$$

It is easy to see that $\rho_{ABR} < 1$ for all choices of η , ϵ , and h . The largest value of ρ_{ABR} occurs for the lowest frequency, where

$$\rho_{ABR} = c \left(\frac{h^2 \eta}{(1-c)(2+hc) + h^2 \eta} \right) \left| \frac{h^2 \eta}{(1-c)(2+hc) + h^2 \eta + h^2 \eta s} \right|. \quad (46)$$

From (39) and some algebraic manipulation, we arrive at an estimate for the spectral radius of the ABR-postconditioned system

$$p = \frac{h^2}{2(1-c)} \left(1 + \frac{h^2}{2} \right) + \left(\frac{h}{(1-c) + is} \right) \epsilon. \quad (45)$$

$$e = \frac{h}{\eta is}, \quad (44)$$

$$b = c, \quad (43)$$

$$a = \frac{h^2}{2(1-c)} \left(1 + \frac{h^2}{2} \right) + \eta, \quad (42)$$

g in (38), with

Using the Fourier transform, it suffices to consider the 2×2 iteration matrix, where $c = \cos(k\pi h)$ and $s = \sin(k\pi h)$.

• ρ_{ABF} in (46) for the ABR-postconditioned matrix, AD^{-1} , in (35);

summarized below:

characterized the spectral radii of the various block Gauss-Seidel iterations as n , and p while (51) is motivated by the quasi-Fermi variant, (8)-(10). We have was motivated by the drift-diffusion equations, (3)-(5), written in terms of n , when h is a fixed positive quantity. As we just noted, the ABR postconditioner In practical problems, it is more interesting to compare the various methods Gauss-Seidel to the ABR-postconditioned system, (35), in the limit as $h \rightarrow 0$. the block Gauss-Seidel procedure applied to (51) is equivalent to applying block Moreover, the limit of ρ_{GF} as $h \rightarrow 0$ is the expression as shown in (49). Hence, It is easy to see that $\rho_{GF} < 1$ and that $\rho_{ABF} < \rho_{GF}$ where ρ_{ABF} is given by (46).

$$\rho_{GF} = \left(\frac{\eta h^2}{2(1-c) + \eta h^2} \right) \left| \frac{(1-c)(2+ch) + chs}{(1-c)ch + chs} \right|. \quad (52)$$

If we postcondition B by ABR, then we can estimate the spectral radius by studying the eigenvalues of the matrix, G , appearing in (36) again; this leads to the same approximation, ρ_{ABF} in (46), for the spectral radius. This follows from ABR's insensitivity to the scaling in (51). (Nevertheless, the ABR technique can be effective when applied to the quasi-Fermi variant of the drift-diffusion equations, as we will see in § 4.) Obviously, we can apply block Gauss-Seidel directly to B . After using the local Fourier trick again and some algebraic manipulation, we find that an estimate for the spectral radius is

$$B = \begin{bmatrix} T & I & I \\ -\eta T & T + \epsilon S & I \\ T + \eta I & 0 & -\eta I \end{bmatrix} = \begin{bmatrix} T + \eta I & \eta \epsilon S & -\eta(T + \epsilon S) \end{bmatrix}, \quad (51)$$

A for the quasi-Fermi variables is simply that system is represented by (8)-(10). The analogous matrix corresponding to equations can also be written in terms of quasi-Fermi variables, n , v , and w ; in terms of the primitive variables, (3)-(5). As noted in § 1, the drift-diffusion Now A in (30) roughly corresponds to the drift-diffusion equations written (5) can diverge, which is in agreement with the results of [21]. These arguments suggest that Newton-Gauss-Seidel applied to (3)-dominant. In order for $\rho_{PV} < 1$, it is necessary which is monotonically increasing in η . In other words, A must be diagonally that $\epsilon = \delta\eta$ for δ sufficiently large or, in other words, A must be diagonally

$$\rho_{PV} = \left| \frac{h^2 \eta}{(1-c)(2+h\epsilon) + h\epsilon s} \right|, \quad (50)$$

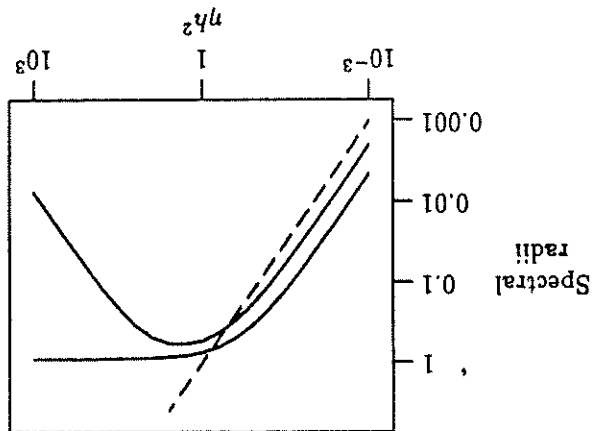
and (39), we arrive at an estimate for the spectral radius

Figure 1 illustrates several important points. First, the ABR approach can slow down for intermediate coupling between the equations, but does well when the equations are weakly or strongly coupled. Our estimates suggest that ABR-Newton-Gauss-Seidel iteration based on quasi-Fermi variables, (51); this superiority is especially dramatic for strongly coupled systems. From this, we surmise that a coupled Newton-ABR-Gauss-Seidel method should converge faster than a simple Gauss-Seidel-Newton method. Nevertheless, we would expect a Newton-Gauss-Seidel approach based on quasi-Fermi variables to converge based on our heuristic analysis, which concurs with the more exact results reported in [15]. Finally, we can see that a Newton-Gauss-Seidel method for the usual system will diverge when the coupling becomes large [21].

Let us emphasize again that $\rho_{ABR} > \rho_{GF}$ which suggests that the ABR approach applied to (3)-(5) or (8)-(10) will converge more rapidly than a Newton-Gauss-Seidel iteration applied to (8)-(10). Figure 1 shows the values of these spectral radii as functions of ηh^2 .

- ρ_{GF} in (52) for the unconditioned quasi-Fermi matrix, B , in (51).
- ρ_{PV} in (50) for the unconditioned matrix, A , in (30);

Figure 1: This figure shows: (1) the ABR spectral radius, $\rho_{ABR}(\eta h^2)$ from (46), as the lower solid line; (2) the block Gauss-Seidel spectral radius, $\rho_{PV}(\eta h^2)$ from (50), as the dashed line; and (3) the block Gauss-Seidel spectral radius for the equations in the quasi-Fermi variables, $\rho_{GF}(\eta h^2)$ from (52), as the upper solid line. Here $c = 0.9$ and $\epsilon h = 2$.



where μ is an eigenvalue of M . Equation (58) holds when $\lambda = 0$ or $\lambda = \alpha^2/(1 + \mu)^2$. For small α , everything is fine; for large α , things can go awry since μ may be quite small.

$$(58) \quad \det \begin{bmatrix} (1 + \mu)\lambda & \alpha\lambda \\ \alpha & (1 + \mu)\lambda \end{bmatrix} = 0,$$

We can also compute the eigenvalues of the iteration matrix for block Gauss-Seidel applied to W in (53) by solving problems of the form

preconditioners, ABR is not a panacea. hold for negative μ , but things are a bit worse in that case. Hence, like other singularities occur and the results then correspond to figure 1. Similar results If you change the sign in one of the off-diagonal blocks to $-\alpha I$ in (53), then no behavior is not characteristic of matrices obtained by discretizing elliptic PDEs. This is not at all surprising since W is also singular for $\alpha = \sqrt{1 + \mu}$; such $\alpha \leq 1 + \mu$ and there is an infinite eigenvalue associated with $\alpha = \sqrt{1 + \mu}$. However, things can break down for $1 \leq 0 < \alpha < 1 + \mu < \infty$. Then, from (57), $|\lambda(\alpha)| < 1$ Let μ be a fixed positive eigenvalue of M . Then, from (57), $|\lambda(\alpha)| < 1$

$$(57) \quad \lambda = \frac{\alpha^2 \mu^2}{1 - \alpha^2 + \mu^2}.$$

Now, (56) holds when $\lambda = 0$ and

$$(56) \quad \det \begin{bmatrix} (1 - \alpha^2 + \mu)\lambda & -\alpha\mu\lambda \\ \alpha\mu & (1 - \alpha^2 + \mu)\lambda \end{bmatrix} = 0.$$

The eigenvalues of the block Gauss-Seidel iteration matrix associated with $W D^{-1}$ in (55) are given by the solutions of the scalar problem

$$(55) \quad W D^{-1} = \begin{pmatrix} 1 \\ 1 - \alpha^2 \end{pmatrix} \begin{bmatrix} (1 - \alpha^2)I + M & -\alpha M \\ -\alpha M & (1 - \alpha^2)I + M \end{bmatrix}.$$

and so

$$(54) \quad D^{-1} = \begin{pmatrix} 1 \\ 1 - \alpha^2 \end{pmatrix} \begin{bmatrix} I & -\alpha I \\ -\alpha I & I \end{bmatrix}$$

of its rows. The ABR postconditioner would then be dimensional grid; in this case, M would have 4 nonzero entries of $-1/4$ in most obtained from discretizing $-\Delta$ by scaled finite differences on a uniform two- in $(-1, 1)$. One possibility for $I + M$ is the block tridiagonal matrix that is Here, M is symmetric with zeros on its diagonal and having eigenvalues, $\{\mu_j\}$,

$$(53) \quad W \doteq \begin{bmatrix} I + M & \alpha I \\ \alpha I & I + M \end{bmatrix} \in \mathbb{R}^{2\nu \times 2\nu}.$$

Let us turn to a simple example that shows how the ABR technique can break down. Consider the block matrix

4 Numerical Experiments

Let us now consider the results of computational experiments on some 'realistic' semiconductor modeling problems. Excluding oxide regions, the PDES actually used in our computer experiments are the drift-diffusion equations based on either the primitive or quasi-Fermi variables

$$(59) \quad -\Delta^2 n + n - p - N = -\Delta^2 n + e^{-n-v} - e^{-w-n} - N = 0,$$

$$(60) \quad \Delta \cdot J_n = 0,$$

$$(61) \quad \Delta \cdot J_p = 0$$

where the current densities are now

$$(62) \quad J_n = \mu_n (n \Delta n - \Delta n) = \mu_n e^{-n-v} \Delta v,$$

$$(63) \quad J_p = -\mu_p (p \Delta p + \Delta p) = -\mu_p e^{-w-n} \Delta w.$$

Here $N(x)$ represents the (net) impurity concentration while μ_n and μ_p are carrier mobility functions. Recall from § 1 that the primitive variables are n, n, n, n and p ; the quasi-Fermi variables are n, v, v and w, w , defined by (6) and (7).

We have performed experiments with several semiconductor structures using both coupled and plug-in methods. The plug-in implementation uses Newton's method, (12) and (13), to solve the nonlinear Poisson equation, based on the Slotboom variables [24], for n . The linear electron continuity equation, (60) and (62), is then solved for n ; p is then determined from the corresponding linear continuity equation. The coupled algorithm is based on applying Newton's method to (59)-(63) in order to obtain either the primitive variables or the quasi-Fermi variables. The coupled-ABF procedure uses block Gauss-Seidel to solve the Newton correction equations, (12). In our experiments, sparse direct methods were always used for the innermost linear equations. We have presented our most recent experience in more detail elsewhere [3, 2].

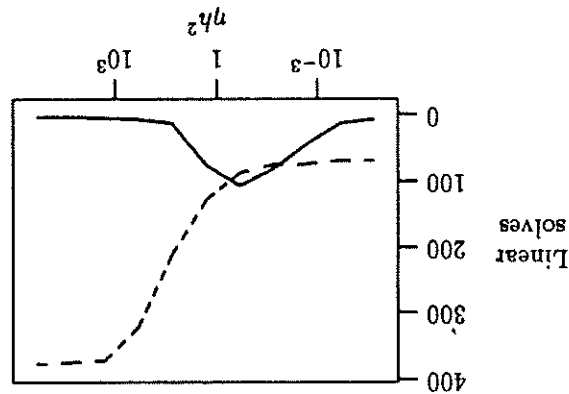
It is difficult to relate the general nonlinear semiconductor problem to linear systems of equations involving the matrix, A , in (30). The mobilities, μ_n , typically are functions of the spatial variable, x , or are nonlinear. The semiconductor equations are normally discretized by an exponentially upwind scheme (for example, see [6, 23]) that effectively reduces the convective terms in the carrier equations, (60)-(63). Ignoring the variable mobilities and meshes, the semiconductor problem discretized with such a scheme gives rise to matrices analogous to A appearing in (30) or B appearing in (51). In the notation of § 3, ηh^2 can be as large as 10^3 or more while $\epsilon h = O(1)$ for realistic semiconductor simulations. These problems are characterized by the extreme variations in the dependent variables and the nonlinearity of the PDES.

Figure 2 shows the results of using plug-in and coupled-ABF for a one-dimensional, one-carrier, resistive-bar problem with a constant mobility. For this problem, the coupled algorithm is based on solving for n and n with the

We repeated the computations for a two-dimensional two-carrier resistive slab with a modest voltage applied across it. In this case, we solved for u , v , and w in (59)-(63) when applying the coupled Newton-ABF method, but retained the same strategy as before for the plug-in algorithm. The nonlinear iterations were terminated when approximately four-digit accuracy was obtained; the inner block Gauss-Seidel iteration was terminated based on an adaptive Newton-Richardson strategy [5, 3]. For successive constant dopings ranging from approximately 10^4 up to 10^8 (in the scaled units), the coupled-ABF algorithm used 6 iterations; moreover, only one or two block Gauss-Seidel iterations were usually required per outer Newton iteration. The plug-in iteration used roughly 50 iterations to obtain the same accuracy for that range of doping values. In [3], we reported results obtained by simulating a large essentially uniformly doped structure in a magnetic field; the results there showed that the ABF approach is quite competitive for these drift-dominated devices.

Newton-ABF-Gauss-Seidel approach. The plug-in and coupled methods were iterated until approximately six-digit accuracy (in the 2-norm) was obtained; for each outer coupled Newton iteration, the ABF-postconditioned block Gauss-Seidel method was iterated until six-digit accuracy was obtained. (In practice, an adaptive Newton-Richardson strategy should be used to terminate the inner Gauss-Seidel iteration [5, 3].) As the answer is approached, the matrices that arise are exactly of the form found in (30). Note that plug-in slows down with increasing coupling and the ABF approach slows for intermediate coupling, which supports the results of § 3 and mimics figure 1.

Figure 2: This plot contrasts the number of linear solves for the plug-in algorithm (dashed line) with the number of linear solves for the Newton-ABF-Gauss-Seidel algorithm (solid line) for a one-dimensional resistive bar. For this idealized problem, η^2 is a measure of the off-diagonal coupling (see (30)).



In summary, the ABF method automatically provides an alternative set of dimensions, scaled grid-function unknowns for use in either a plug-in iteration or in the inner iteration of a coupled Newton scheme. These new grid functions are found by examining the coupling between the original unknown grid functions locally at each grid point.

For either approach (plug-in or coupled), the selection of new variables (that is, the computation of the postconditioner, D^{-1}) can be carried out at every

5 Conclusions

In table 1, some results from the simulation of a small two-dimensional MOS transistor in a high-current state are summarized. We made use of the same solution strategies as just described for the two-dimensional resistive slab. For the coarser grid consisting of 163 points, the coupled-ABF method reduced the number of linear solves as compared with the plug-in method. The coupled approach used fewer iterations, but required the solution of a matrix of order 3 times larger than the plug-in and coupled-ABF methods. For the finer grid, the ABF scheme was able to substantially reduce both the nonlinear and linear work as compared to the plug-in method; we do not expect such spectacular improvements as the mesh spacing becomes smaller, but it is clear that the finer mesh in this case enabled the ABF postconditioner to almost decouple the system. Once again, the ABF preconditioner was remarkably effective for this drift-dominated problem, and only one or two block Gauss-Seidel iterations were usually required per outer Newton iteration.

We also simulated a forward-biased pn junction with low doping, where diffusion effects are more prominent. The coupled-ABF approach has more difficulty with this problem. It is sometimes necessary to do 12 or 15 block Gauss-Seidel iterations for each Newton iteration. Further experiments on pn junctions and bipolar transistors are needed.

Nonlinear algorithm	Vertices	Nonlinear iterations	Linear iterations
coupled	163	14	-
plug-in	163	42	217
coupled-ABF	163	41	200
	2765	41	123
	2765	26	78

Table 1: Results from various algorithms for a small MOS transistor in saturation.

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iteration step, or less frequently, to the extreme of being computed once at the beginning and then held fixed throughout the calculation. Indeed, as plug-in and coupled are generic terms covering a wide classes of potential algorithms, the number of possible APF schemes is really unbounded. In fact, the obvious nonlinear variant or just doing plug-in on the alternately blocked system may prove to be attractive.

The heuristic analysis of some model problems in § 3 (born out by the results in § 4) suggests that the APF preconditioner avoids some of the convergence difficulties associated with plug-in algorithms for tightly coupled PDEs. With some exceptions, our results also suggest that a Newton-Gauss-Seidel iteration applied to the drift-diffusion equations written in terms of the primitive variables should converge faster if the matrices are preconditioned by APF; the exceptions occur for cases of moderate coupling when diffusion plays a more significant role. We also considered the drift-diffusion equations written in quasi-Fermi variables where it appears that APF preconditioning should always improve the convergence rate of a Newton-Gauss-Seidel iteration. For drift-diffusion modeling, our analyses further validate the usual engineering practice of performing plug-in on a nonlinear version of the drift-diffusion equations (also see [14]).

The APF preconditioner shares some similarities with 'element-by-element' preconditioners [10, 11, 25] and the recently introduced class of 'transforming smoothers' [26]. The APF technique is not identical to the 'element-by-element' approach, as far as we can determine, but is a special instance of r -transforming smoothers. We have essentially ignored incomplete preconditioners [18, 1, 17] since we are interested in computations where it is difficult to store just the Jacobian and a few solution vectors in main memory. It has also been suggested that a nonlinear GMRES algorithm can be used to accelerate the plug-in algorithm [16]. Further experimentation and comparisons are needed.

The decisions as to what overall approach to follow, what particular inner and outer iterations to use, and how to incorporate the APF technique all seem to be highly problem dependent, and the usual array of empirical trade-offs must be taken into account. What we want to emphasize here is that APF is a simple and easily implemented preconditioner that can have a dramatic effect on the convergence rate of commonly used plug-in and coupled algorithms for systems of nonlinear elliptic PDEs.

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